5 Finite-Length Discrete Transform

Introduction

- In this chapter, *finite-length transforms* are discussed
- In practice, it is often convenient to map a finitelength sequence from time domain into a finitelength sequence of the same length in the frequency domain
- The samples of the forward transform are unique and represented as a linear combination of the samples of the time domain sequence
- The samples of the inverse transform are obtained similarly from the samples of the transform domain

Introduction

- In some applications, a very long-length time domain sequence is broken up into a set of shortlength sequences and a finite-length transform is applied to each short-length sequence
- The transformed sequences are processed in the transform domain
- Time domain equivalents are produced using the inverse transform
- The processed short-length sequences are grouped together in the time domain to form the final long-length sequence

- Let *x*[*n*] denote a length-*N* time domain sequence with *X*[*k*] denoting the coefficients of its *N*-point orthogonal transform
- A general form of the orthogonal transform pair is of the form

$$X[k] = \sum_{n=0}^{N-1} x[n] \psi^*[k,n], \quad 0 \le k \le N-1$$

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] \psi[k,n], \quad 0 \le n \le N-1$$

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- In the transform pair, the *basis sequences ψ*[*k*,*n*] are also length-*N* sequences in both domains
- In the class of finite-dimensional transforms, the basis sequences satisfy the condition

$$\frac{1}{N}\sum_{n=0}^{N-1}\psi[k,n]\psi^*[l,n] = \begin{cases} 1, & l=k\\ 0, & l\neq k \end{cases}$$

Basis sequences ψ[k,n] satisfying the above condition are said to be *orthogonal* to each other

- An important consequence of the orthogonality of the basis sequence is the energy preservation property of the transform
- The energy $\sum_{n=0}^{N-1} |x[n]|^2$ of the time domain sequence x[n] can be computed in the transform domain
- The energy can be written as

$$\sum_{n=0}^{N-1} |x[n]|^2 = \sum_{n=0}^{N-1} x[n] x^*[n]$$

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• Let us express *x*[*n*] in terms of its transform domain representation

$$\sum_{n=0}^{N-1} x[n] \ x^*[n] = \sum_{n=0}^{N-1} \left(\frac{1}{N} \sum_{k=0}^{N-1} X[k] \psi[k,n] \right) \ x^*[n]$$

$$= \frac{1}{N} \sum_{k=0}^{N-1} X\left[k \left(\sum_{n=0}^{N-1} x^*[n] \psi[k,n]\right) = \frac{1}{N} \sum_{k=0}^{N-1} X[k] X^*[k]\right]$$

$\sum_{n=0}^{N-1} x[n] ^2$	$=\frac{1}{N}\sum_{k=0}^{N-1} \left X[k] \right ^2$
n=0	k=0

which is known as the Parseval's relation

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The Discrete Fourier Transform

- In the following, the discrete Fourier transform, DFT, is defined
- The inverse transformation, IDFT is developed
- Some important properties of the DFT are discussed
- DFT has several important applications:
 - Numerical calculation of the Fourier transform in an efficient way
 - Implementation of linear convolution using finite-length sequences

Definition of the DFT

 Discrete Fourier transform (DFT) of the length-N sequence x[n] is defined by

$$X[k] = \sum_{n=0}^{N-1} x[n] e^{-j2\pi kn/N}, \quad 0 \le k \le N-1$$

- The basis sequences are: $\psi[k,n] = e^{-j2\pi kn/N}$ which are complex exponential sequences
- As a result, DFT coefficients *X*[*k*] are complex numbers, even if *x*[*n*] are real
- It can be easily shown that the basis sequences $e^{j2\pi kn/N}$ are orthogonal

Discrete Fourier Transform (DFT)

- Common notation with DFT: $W_N = e^{-j2\pi/N}$
- DFT can now be written as follows:

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn}, \quad k = 0, 1, ..., N-1$$

Inverse Discrete Fourier Transform (IDFT):

$$x[n] = \frac{1}{N} \sum_{k=0}^{N-1} X[k] W_N^{-kn}, \quad n = 0, 1, \dots, N-1$$

• *X*[*k*] and *x*[*n*] are both sequences of finite-length *N*

Discrete Fourier Transform Pair

• The analysis equation:

$$X[k] = \sum_{n=0}^{N-1} x[n] W_N^{kn}, \quad k = 0, 1, ..., N-1$$

• The synthesis equation:

$$x[n] = \frac{1}{N} \sum_{n=0}^{N-1} X[k] W_N^{-kn}, \quad n = 0, 1, ..., N-1$$

DFT

• DFT pair is denoted as: $x[n] \leftrightarrow X[k]$

Relation Between the Discrete-Time Fourier Transform and the DFT

 The DTFT X(e^{j∞}) of the length-N sequence x[n] defined for 0≤n≤N-1 is given by:

$$X(e^{j\omega}) = \sum_{n=-\infty}^{\infty} x[n] e^{-j\omega n} = \sum_{n=0}^{N-1} x[n] e^{-j\omega n}$$

• By uniformly sampling $X(e^{j\omega})$ at N equally spaced frequencies $\omega_k = 2\pi k/N$, $0 \le k \le N-1$, on the ω -axis between $0 \le k \le 2\pi$

$$X(e^{j\omega})\Big|_{\omega=2\pi k/N} = \sum_{n=-\infty}^{\infty} x[n] e^{-j2\pi nk/N}, \quad 0 \le k \le N-1$$

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Relation Between the Discrete-Time Fourier Transform and the DFT

- The *N*-point DFT sequence X[k] is precisely the set of frequency samples of the Fourier transform $X(e^{j\omega})$ of the length-*N* sequence x[n] at *N* equally spaced frequencies, $\omega_k = 2\pi k/N$, $0 \le k \le N-1$
- Hence, the DFT *X*[*k*] represents a frequency domain representation of the sequence *x*[*n*]
- Since the computation of the DFT samples involve a finite sum, for time domain sequences with finite sample values, the DFT always exists

Numerical Computation of the Fourier Transform Using the DFT

- The DFT provides a practical approach to the numerical computation of the Fourier transform of a finite-length sequence
- Let X(e^{jw}) be the Fourier transform of a length-N sequence x[n]
- We wish to evaluate $X(e^{j\omega})$ at a dense grid of frequencies $\omega_k = 2\pi k/M$, $0 \le k \le M-1$, where M >> N:

$$X(e^{j\omega}) = \sum_{n=-\infty}^{\infty} x[n] \ e^{-j\omega_k n} = \sum_{n=0}^{N-1} x[n] \ e^{-j2\pi kn/M}$$

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Numerical Computation of the Fourier Transform Using the DFT

• Define a new sequence $x_e[n]$ obtained from x[n] by augmenting with *M*-*N* zero-valued samples

$$x_e[n] = \begin{cases} x[n], & 0 \le n \le N-1 \\ 0, & N \le n \le M-1 \end{cases}$$

• Making use of $x_e[n]$ we obtain

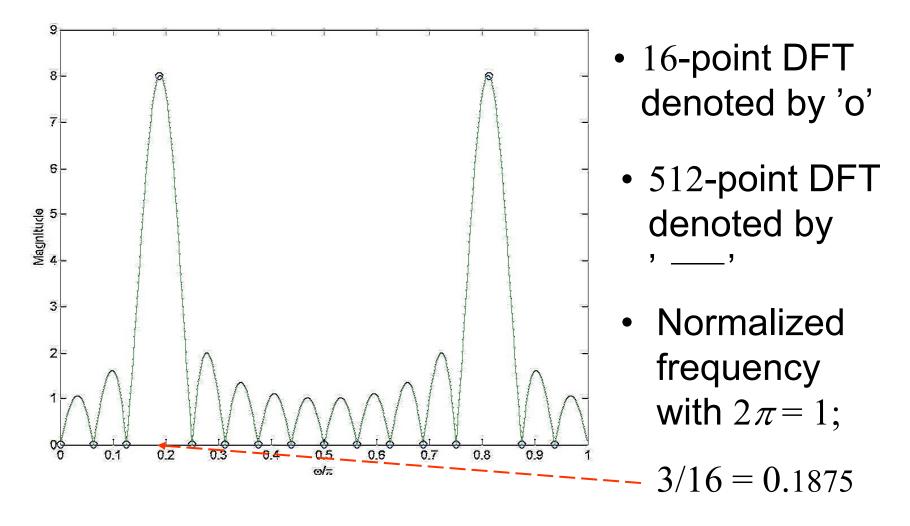
$$X(e^{j\omega_{k}}) = \sum_{n=0}^{M-1} x_{e}[n] e^{-j2\pi kn/M}$$

which is an *M*-point DFT $X_e[k]$ of the length-*M* sequence $x_e[n]$

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Example 5.5:

- Compute the *N*-point DFT of the length-16 sequence $x[n]=\cos(6\pi n/16)$ of angular frequency $\omega_0=0.375\pi$
- The 16-point DFT of x[n] is (Example 5.2) $X[k] = \begin{cases} 8, & \text{for } k = 3 \text{ and } k = 13 \\ 0, & \text{otherwise} \end{cases}$
- Since the Fourier transform X(e^{jω}) is a continuous function of ω, we can plot it more accurately by computing the DFT of the sequence x[n] at a dense grid of frequencies using MATLAB



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Sampling the Fourier Transform

• The discrete Fourier transform DFT can also be obtained by sampling the discrete-time Fourier transform DTFT, $X(e^{j\omega})$, uniformly on the ω -axis between $0 \le \omega \le 2\pi$, at $\omega_k = 2\pi k/N$, $k=0,1,\ldots,N-1$

$$X[k] = X(e^{j\omega})\Big|_{\omega = 2\pi k/N}$$

= $\sum_{n=0}^{N-1} x[n]e^{-j2\pi kn/N}$, $k = 0, 1, ..., N-1$

• *X*[*k*] is now a finite-length sequence of length *N* like the time domain sequence *x*[*n*]

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Operations on Finite-Length Sequences

- Like the Fourier transform, the DFT also satisfies a number of properties that are useful in signal processing
- Some of the properties are essentially identical to those of the Fourier transform, while some others are different
- Differences between two important properties are discussed:
 - Shifting and
 - Convolution

- Several DFT properties and theorems involve shifting in the time domain and in the frequency domain
- The operation of shifting of a finite-length sequence in time domain is referred to as *circular time-shifting*
- In frequency domain the corresponding operation is referred to as *circular frequency-shifting*

- Consider length-N sequences defined for the range 0≤n≤N-1
- Such sequences have zero for n < 0 and $n \ge N$
- Shifting such a sequence x[n] for any arbitrary integer n_0 , the resulting sequence $x_1[n]=x[n-n_0]$ is no longer defined for the range $0 \le n \le N-1$
- It is necessary to define a shift operation that will keep the shifted sequence in the range 0≤n≤N-1
- This is achieved using the *modulo operation*

- Let 0,1, ..., *N* be a set of *N* positive integers and let *m* be any integer
- The integer *r* obtained by evaluating *m* modulo *N* is called the *residue* and it is an integer with a value between 0 and *N*-1
- The modulo operation is denoted as

 $\langle m \rangle_N = m \text{ modulo } N$

• If we let $r = \langle m \rangle_N$, then r = m + lNwhere *l* is an integer to make m+lN a number in the range $0 \le n \le N-1$

 Using the modulo operation, the circular shift of a length-N sequence x[n] can be defined using the equation

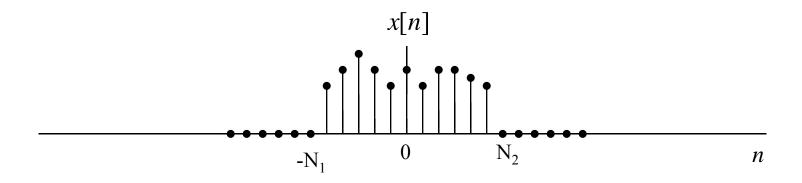
$$x_{C}[n] = x[\langle n - n_{0} \rangle_{N}]$$

where *x*[*n*] is also length-*N* sequence

 The concept of circular shift of a finite-length sequence corresponds to "rotation" of the sequence within the interval 0≤n≤N-1

Representation of a Finite-Length Sequence

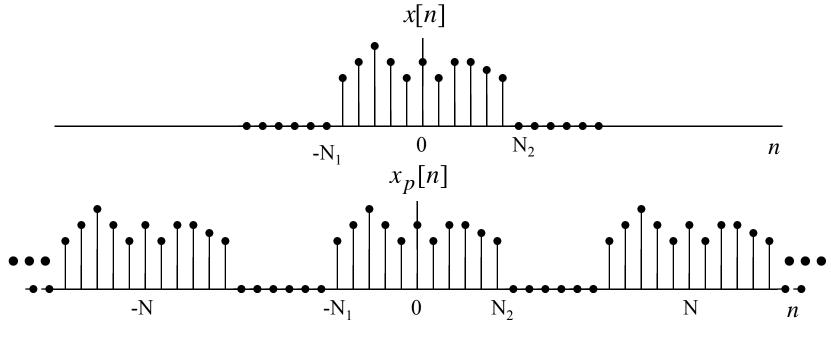
• Consider a general sequence x[n] that is of *finite-length*, i.e., for some integers N_1 and N_2 , x[n] = 0 outside the range $-N_1 \le n \le N_2$



• The shifting operation of finite-length sequences can be represented via periodic sequences

Representation of Aperiodic Signals

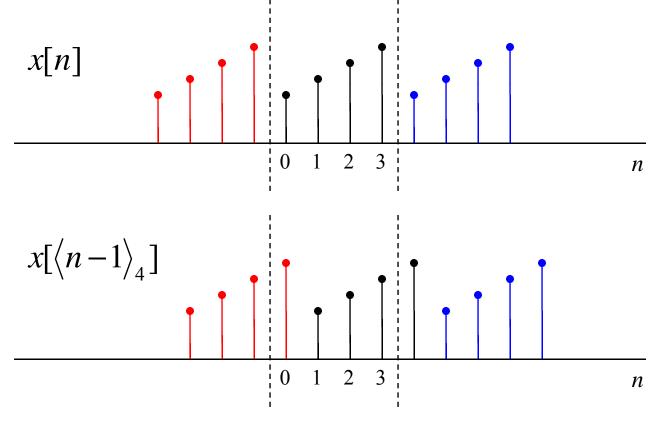
A periodic sequence, x_p[n], is formed from the *aperiodic* sequence with x[n] as one period



As N approaches infinity, x_p[n] = x[n] for any finite value n

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Shifting of a finite sequence corresponds to rotation



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Circular Convolution

- Consider two length-*N* sequences, *g*[*n*] and *h*[*n*]
- Their linear convolution is a sequence of length 2*N*-1

$$y_L[n] = \sum_{m=0}^{N-1} g[m]h[n-m], \quad n = 0, 1, ..., 2N-1$$

• In order to calculate the above linear convolution both length-*N* sequences have been zero-padded to extend their length to 2*N*-1

Circular Convolution

A convolution-like operation resulting in a length-N sequence y_C[n], called a *circular convolution* is defined as

$$y_C[n] = \sum_{m=0}^{N-1} g[m]h[\langle n-m \rangle_N]$$

- The above operation is often referred to as an *N-point circular convolution*
- Due to length-*N* sequences, the *N*-point circular convolution is denoted as

$$y_C[n] = g[n] \otimes h[n] = h[n] \otimes g[n]$$

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Application of Circular Convolution

- The *N-point circular convolution* does not correspond to the *linear convolution* of two length-*N* sequences
- The circular convolution can, however, be used to compute the linear convolution correctly:
 - The linear convolution of two finite-length sequences of length N and M results in a sequence of length N+M-1
 - The circular convolution must be computed for the length *N*+*M*-1 by zero-padding the original sequences

Classification of Finite-Length Sequences

- For a finite-length sequence defined for 0≤n≤N-1, all definitions of symmetry do not apply
- The definitions of symmetry in the case of finitelength sequences are given such that the symmetric and antisymmetric parts of length-*N* sequence are also of length *N* and defined for the same range of values of the time index *n*

Classification Based on Geometric Symmetry

- Geometric symmetry is an important property in DSP, i.e., in the properties of FIR filters
- A length-*N* symmetric sequence *x*[*n*] satisfies the condition

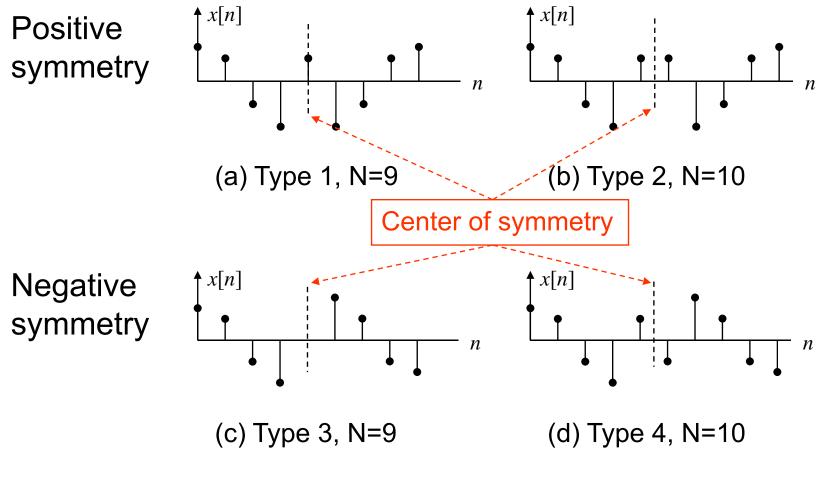
$$x[n] = x[N-1-n]$$

• A length-*N* antisymmetric sequence *x*[*n*] satisfies the condition

$$x[n] = -x[N-1-n]$$

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Geometric Symmetry of Sequences



Type 1 Symmetry with Odd Length

- Type 1 symmetric sequence, with N=9, is x[n] = x[0] + x[1] + x[2] + x[3] + x[4] + x[5] + x[6] + x[7] + x[8]
- The Fourier transform is

$$X(e^{j\omega}) = x[0] + x[1]e^{-j\omega} + x[2]e^{-j2\omega} + x[3]e^{-j3\omega} + x[4]e^{-j4\omega} + x[5]e^{-j5\omega} + x[6]e^{-j6\omega} + x[7]e^{-j7\omega} + x[8]e^{-j8\omega}$$

• Now, x[0]=x[8], x[1]=x[7], x[2]=x[6], x[3]=x[5]

$$X(e^{j\omega}) = x[0](1 + e^{-j8\omega}) + x[1](e^{-j\omega} + e^{-j7\omega}) + x[2](e^{-j2\omega} + e^{-j6\omega}) + x[3](e^{-j3\omega} + e^{-j5\omega}) + x[4]e^{-j4\omega}$$

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Type 1: Symmetry with Odd Length

 Taking e^{-j4w} as a common factor in each group of terms

$$\begin{split} X(e^{j\omega}) &= x[0]e^{-j4\omega} \left(e^{j4\omega} + e^{-j4\omega} \right) + x[1]e^{-j4\omega} \left(e^{-j3\omega} + e^{-j3\omega} \right) \\ &+ x[2]e^{-j4\omega} \left(e^{j2\omega} + e^{-j2\omega} \right) + x[3]e^{-j4\omega} \left(e^{j\omega} + e^{-j\omega} \right) + x[4]e^{-j4\omega} \\ X(e^{j\omega}) &= e^{-j4\omega} \left\{ x[0] \left(e^{j4\omega} + e^{-j4\omega} \right) + x[1] \left(e^{-j3\omega} + e^{-j3\omega} \right) \\ &+ x[2] \left(e^{j2\omega} + e^{-j2\omega} \right) + x[3] \left(e^{j\omega} + e^{-j\omega} \right) + x[4] \right\} \end{split}$$

$$X(e^{j\omega}) = e^{-j4\omega} \{2x[0]\cos(4\omega) + 2x[1]\cos(3\omega) + 2x[2]\cos(2\omega) + 2x[3]\cos(\omega) + x[4]\}$$

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Type 1: Symmetry with Odd Length

- Notice that the quantity inside the braces, { }, is a real function of ω and can assume positive or negative values in the range $0 \le \omega \le \pi$
- The of the sequence is given by $\theta(\omega) = -4\omega + \beta$ where β is either 0 or π , and hence the phase is a linear function of ω
- In general, for Type 1 linear-phase sequence of length-*N*

$$X(e^{j\omega}) = e^{-j(N-1)\omega/2} \left\{ x\left[\frac{N-1}{2}\right] + 2\sum_{n=1}^{(N-1)/2} x\left[\frac{N-1}{2} - n\right] \cos(\omega n) \right\}$$

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Type 2: Symmetry with Even Length

• Similarly, the Fourier transform of Type 2 symmetric sequence, with *N*=8, can be written

$$X(e^{j\omega}) = e^{-j7\omega/2} \left\{ 2x[0]\cos(\frac{7\omega}{2}) + 2x[1]\cos(\frac{5\omega}{2}) + 2x[2]\cos(\frac{3\omega}{2}) + 2x[3]\cos(\frac{\omega}{2}) \right\}$$

where the phase is given by $\theta(\omega) = -\frac{7\omega}{2} + \beta$

• In general, for Type 2 linear-phase sequence of length-*N*

$$X\left(e^{j\omega}\right) = e^{-j(N-1)\omega/2} \left\{ 2\sum_{n=1}^{N/2} x\left[\frac{N}{2} - n\right] \cos\left(\omega(n-\frac{1}{2})\right) \right\}$$

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Type 3: Antisymmetry with Odd Length

- The Fourier transform of Type 3 antisymmetric sequence, with N=9, is (notice that x[4]=0) $X(e^{j\omega}) = x[0] + x[1]e^{-j\omega} + x[2]e^{-j2\omega} + x[3]e^{-j3\omega} + x[4]e^{-j4\omega} + x[5]e^{-j5\omega} + x[6]e^{-j6\omega} + x[7]e^{-j7\omega} + x[8]e^{-j8\omega}$
- Now, x[0]=-x[8], x[1]=-x[7], x[2]=-x[6], x[3]=-x[5] and x[4]=0 $X(e^{j\omega}) = x[0](1-e^{-j8\omega}) + x[1](e^{-j\omega}-e^{-j7\omega})$ $+ x[2](e^{-j2\omega}-e^{-j6\omega}) + x[3](e^{-j3\omega}-e^{-j5\omega})$ $X(e^{j\omega}) = e^{-j4\omega} \{x[0](e^{j4\omega}-e^{-j4\omega}) + x[1](e^{j3\omega}-e^{-j3\omega})$ $+ x[2](e^{j2\omega}-e^{-j2\omega}) + x[3](e^{j\omega}-e^{-j\omega})\}$

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Type 3: Antisymmetry with Odd Length

• Multiplying by $j=e^{j\pi/2}$ and 2, we obtain

$$X(e^{j\omega}) = e^{-j4\omega} e^{j\pi/2} \left\{ 2x[0]_{\frac{1}{2j}} \left(e^{j4\omega} - e^{-j4\omega} \right) + 2x[1]_{\frac{1}{2j}} \left(e^{j3\omega} - e^{-j3\omega} \right) \right. \\ \left. + 2x[2]_{\frac{1}{2j}} \left(e^{j2\omega} - e^{-j2\omega} \right) + 2x[3]_{\frac{1}{2j}} \left(e^{j\omega} - e^{-j\omega} \right) \right\}$$

which results in

$$X(e^{j\omega}) = e^{j(-4\omega + \pi/2)} \{ 2x[0]\sin(4\omega) + 2x[1]\sin(3\omega) + 2x[2]\sin(2\omega) + 2x[3]\sin(\omega) \}$$

The phase is now $\theta(\omega) = -4\omega + \frac{\pi}{2} + \beta$

• The antisymmetry introduces a phase shift of $\pi/2$

Type 3 and 4: Antisymmetry with Odd and Even Length

• In general, the Fourier transform of Type 3 linearphase antisymmetric sequence of odd length-*N* is

$$X\left(e^{j\omega}\right) = je^{-j(N-1)\omega/2} \left\{ 2\sum_{n=1}^{(N-1)/2} x\left[\frac{N-1}{2} - n\right] \sin(\omega n) \right\}$$

• Similarly, the Fourier transform of Type 4 linearphase antisymmetric sequence of even length-*N* is

$$X\left(e^{j\omega}\right) = je^{-j(N-1)\omega/2} \left\{ 2\sum_{n=1}^{N/2} x\left[\frac{N}{2} - n\right] \sin\left(\omega(n-\frac{1}{2})\right) \right\}$$

• In both cases, $j=e^{j\pi/2}$ introduces a phase shift of $\pi/2$

Discrete Fourier Transform Theorems

- The important theorems hold for DFT with time domain sequences length-N and their DFTs of length-N, e.g.,
 - Linearity
 - Circular time-shifting
 - Circular frequency-shifting
 - Circular convolution
 - Modulation
 - Parseval's theorem
- The proofs are straightforward using the definitions

Linear Convolution of Two Finite-Length Sequences

- Let *g*[*n*] and *h*[*n*] be two finite-length sequences of lengths *N* and *M*, respectively
- The objective is to implement their linear convolution

$$y_L[n] = g[n] \circledast h[n]$$

- The length of the sequence $y_L[n]$ is L=N+M-1
- The linear convolution can be obtained using the circular convolution with the correct length equal to *L*

Linear Convolution of Two Finite-Length Sequences

 Define two length-L sequences g_e[n] and h_e[n] by appending g[n] and h[n] with zero-valued samples

$$g_{e}[n] = \begin{cases} g[n], & 0 \le n \le N - 1 \\ 0, & N \le n \le L - 1 \end{cases}$$
$$h_{e}[n] = \begin{cases} h[n], & 0 \le n \le M - 1 \\ 0, & M \le n \le L - 1 \end{cases}$$

• Then,

$$y_L[n] = y_C[n] = g_e[n] \square h_e[n]$$

Linear Convolution of Two Finite-Length Sequences Using the DFT

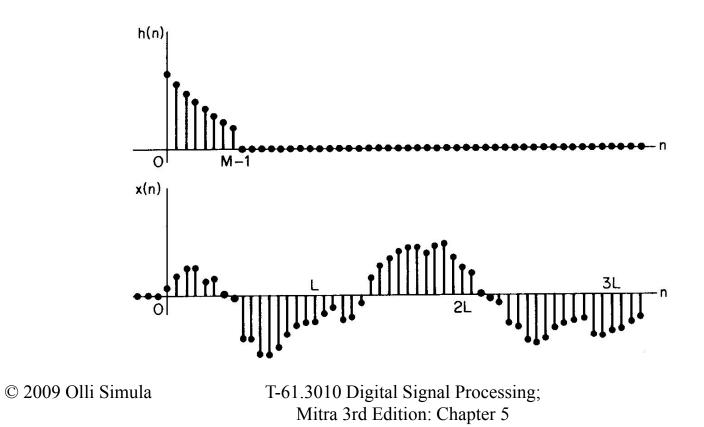
• The linear convolution of two finite-length sequences *g*[*n*] and *h*[*n*] can be implemented using the DFTs of length *L*=*N*+*M*-1 as follows

$$g[n] \quad Zero-padding \quad g_e[n] \quad (N+M-1)-point \text{ DFT} \quad (N+M-1)-y_L[n]$$

$$h[n] \quad Zero-padding \quad h_e[n] \quad (N+M-1)-point \text{ IDFT} \quad point \text{ IDFT} \quad point \text{ IDFT} \quad Length-(N+M-1)$$

Data Sequence of Unknown Length

 Problem: Filtering of a data sequence of unknown, or infinite length with an FIR filter, with impulse response, *h*[*n*], of length *M* using the DFT



Linear Convolution of Finite-Length Sequences

- Filtering of a data sequence of unknown (infinite) length with an FIR filter, with impulse response, *h*[*n*], of length *M* can be implemented via circular convolution, i.e., using the DFT
- The data sequence *x*[*n*] is first segmented into finite-length sections of length-*L*
- Two methods to implement the linear convolution
 - Overlap-add method
 - Overlap-save method

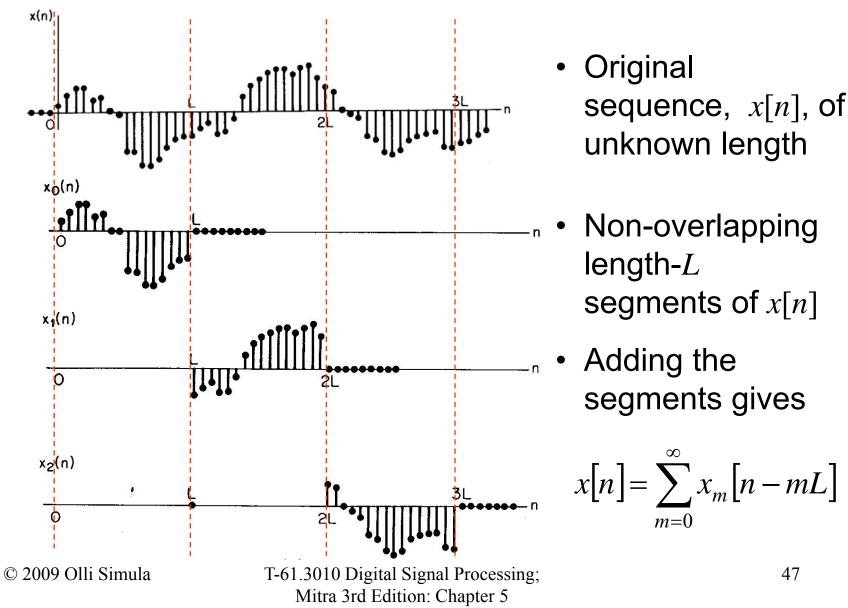
- The causal data sequence *x*[*n*] is first segmented into segments of length *L*
- The original sequence *x*[*n*] can now be written as

$$x[n] = \sum_{m=0}^{\infty} x_m [n - mL]$$

where

$$x_m[n] = \begin{cases} x[n+mL], & 0 \le n \le L-1 \\ 0, & \text{otherwise} \end{cases}$$

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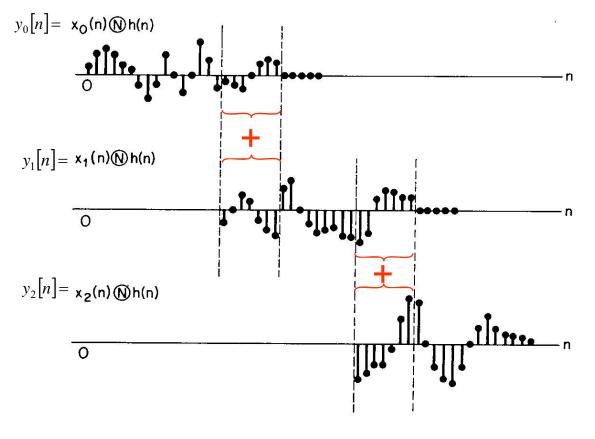


• Substituting the segmented form of *x*[*n*] into the convolution sum

$$y[n] = \sum_{k=0}^{M-1} h[k] x[n-k] = \sum_{k=0}^{M-1} h[k] \sum_{m=0}^{\infty} x_m [n-k-mL]$$
$$= \sum_{m=0}^{\infty} \left(\sum_{k=0}^{M-1} h[k] x_m [n-k-mL] \right) = \sum_{m=0}^{\infty} y_m [n-mL]$$
where $y_m[n] = h[n] \circledast x_m[n]$

• The linear convolutions of *h*[*n*] and the segments of *x_m*[*n*], which all are all of length-*N*, (*N*=*M*+*L*-1) are thus added

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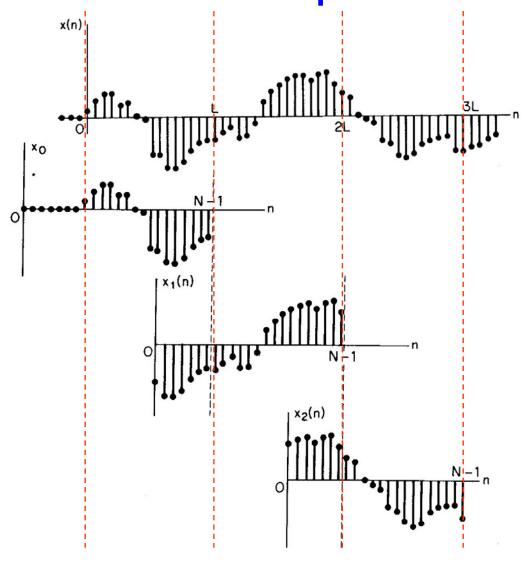
- The linear length-N convolutions of h[n] and x_m[n]
- The overlapping parts of the linear convolutions are added

$$y[n] = \sum_{m=0}^{\infty} y_m[n - mL]$$

Overlap-Save Method

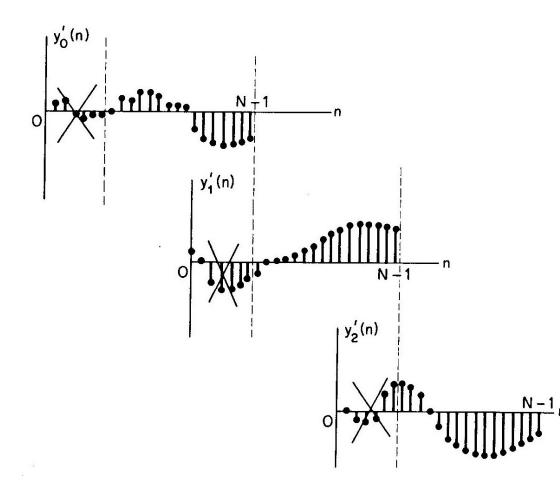
- It is possible to implement the linear convolution also by performing circular convolutions of length shorter than (*M*+*L*-1)
- In this case, it is necessary to segment the original sequence *x*[*n*] is into overlapping blocks *x_m*[*n*],
- The terms of the circular convolution of *h*[*n*] with $x_m[n]$ that correspond to the terms obtained by a linear convolution of *h*[*n*] and $x_m[n]$
- The other, incorrect, terms of the circular convolution are thrown away

Overlap-Save Method



- Original sequence, x[n], of unknown length
- Overlapping length-N segments of x_m[n]
- Circular convolution is implemented with length *N*

Overlap-Save Method



- The length-N circular convolutions of length-M impulse response, h[n], and the blocks x_m[n] of length-N
- The incorrect *M*-1 first terms in each circular convolution are rejected

Summary

- The discrete Fourier transform, DFT, of a finitelength sequence was discussed
- The length of the transform coefficient sequence, i.e., the length of the DFT, is the same as the length of the discrete-time sequence
- The DFT is widely used in a number of digital signal processing applications
- In practice, the DFT can be efficiently implemented using the Fast Fourier Transform (FFT) algorithm